

# Can we Make School Choice more Efficient?

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November 4, 2010

Thousands of students in Boston and New York participate in the public school choice system

- Students (strategically) submit a (strict) preference over schools
- Schools (non-strategically) have (weak) priorities over students
- A central mechanism produces the best *stable* matching for students
  - Individually rational
  - Not blocked: no unmatched student-school pair strictly improves by matching to each other

- Economists recently redesigned the NYC and Boston school admission systems (Abdulkadiroğlu et al. 2005a,b)
- The current design uses the Deferred Acceptance algorithm, with random (single lottery) tie breaking (DA-STB)
  - Strategy-proof for students
  - Stable
  - Inefficient - *About 1,500 students in NYC annually could be assigned to a more preferred school (Abdulkadiroğlu et al. 2009)*

- Student Optimal Stable Matchings (SOSM) can be implemented using the Stable Improvement Cycles mechanism (SIC) (Erdil, Ergin AER2008)
  - Pareto dominates DA-STB
  - *Not strategy-proof for students*
  - Stable (under truth-telling)
  - Efficient (under truth-telling)

*"Nothing is yet known about what kinds of preferences one could expect to be strategically submitted to such a mechanism, or what their welfare consequences would be. Consequently, there is room for more work to further illuminate the tradeoff between efficiency and strategy-proofness."*

*(Abdulkadiroğlu, Pathak, Roth AER2009)*

- Two papers
- Equilibrium behavior of the SIC mechanism
- Approximate large school choice problems using a continuum framework
  - Azevedo-Leshno 2010b develop the theory for college admission with a continuum of students, and show when it approximates large discrete problems

- Using a continuum model, analyze equilibrium of SIC :
  - can *Pareto dominate*, can be *Pareto dominated*, or not be Pareto comparable with DA-STB outcomes
  - assignment be inefficient and unstable with respect to the true preferences
  - simple (intuitive) manipulations, rely only on aggregate information
  - SIC is manipulable in large markets
- Trade changes the valuation
  - students rank schools by their "trade value"
  - results in missallocation
  - trade does not always correct all the missallocation

**DA-STB is inefficient, SOSM is not Strategy-proof:** Erdil, Ergin AER 08;  
Abdulkadiroğlu, Pathak, Roth AER 09

**Current Design:** Abdulkadiroğlu, Pathak, Roth AER PP 2005  
Abdulkadiroğlu, Pathak, Roth, Sönmez AER PP 05

**Incentive properties of deferred acceptance:** Kojima, Pathak AER 09;  
Immorlica, Mahdian SODA 05; Dubins, Freedman AMM 81, Roth MOR  
82

**Experiments:** Che, Sönmez JET 06; Featherstone, Niederle 08; Echenique,  
Wilson, Yariv 09

**Alternative mechanisms:** Abdulkadiroğlu, Che, Yasuda AER forthcoming;  
Kesten QJE forthcoming

**Continuum matching models:** Abdulkadiroğlu, Che, Yasuda 09; Miralles 08



# School Choice (with Continuum of Students)

A school choice problem is:

- Finite set of schools  $S = \{a, b, \dots, c\}$
- Quota  $q_a$  for each school  $a$
- A student (type)  $\theta$  is described by  $\theta = (\prec_\theta, e_\theta)$ 
  - $\prec_\theta$  - a strict preference ordering over  $S \cup \{\phi\}$   
(  $\phi$  denotes being unmatched )
  - $e_\theta$  - vector of priorities  $e_\theta \in \mathbb{R}^{|S|}$   
 $e(a)$  is the priority of the student at school  $a$   
 $\theta$  is acceptable for school  $a \in S$  if  $e_\theta(a) \geq 0$
- A measure  $\eta [m]$  over the set of student types  $\Theta$

# School Choice (with Continuum of Students)

A *matching* is a function  $\mu : \Theta \rightarrow S \cup \{\phi\}$  such that:

- A school  $a$  is matched to the set of students  $\mu^{-1}(a)$   
(a student is matched to a school iff the school is matched to the student)
- School capacity is respected:

$$|\mu(a)| = \eta(\{\theta | \mu(\theta) = a\}) \leq q_a$$

- It is right continuous (technical)

# School Choice (with Continuum of Students)

A pair  $(a, \theta)$  blocks the match  $\mu$  if:

- The student prefers the school over his current match:  $a \succ \mu(\theta)$
- The school is either strictly under capacity  $|\mu(a)| < q_a$  and  $\theta$  is acceptable  $e_\theta(a) \geq 0$

**or**

there is  $\theta'$  such that  $\mu(\theta') = a$  but  $e'(a) < e(a)$

A matching  $\mu$  is *stable* if:

- It is individually rational
- There is no pair  $(a, \theta)$  that blocks

# DA-STB (the current mechanism)

Break ties in priorities and run the Gale-Shapley algorithm with students proposing:

**stage 0** Single random lottery number for each student. Break ties in favor of students with higher lottery number

⇒ Strict priorities for all schools

**stage 1** Students apply to most preferred school that didn't reject them.

**stage 2** Schools tentatively keep the highest priority students up to capacity, and reject the rest

If students were rejected, go to 1 ; otherwise, end.

# The SIC Mechanism

Any efficient mechanism that Pareto dominates DA (SOSM) can be implemented by the SIC algorithm:

- ▶ Run DA-STB
  - ▶ Trade: Find a Pareto improvement that won't violate stability (a stable improvement cycle) and implement it.
- Repeat till none exist

The outcome is an efficient stable matching that Pareto dominates the DA-STB match

# Example (Erdil-Ergin)

School Capacities:  $q_a = q_b = 1, q_z = \infty$

Student preferences:

$\alpha$	$\beta$	$\zeta$
a	b	a
b	a	z
z	z	

Priorities:

	$\alpha$	$\beta$	$\zeta$
a	0	1	0
b	1	0	0
z	0	0	0

## Example (Erdil-Ergin): DA-STB

The only tie break that matters is between  $\alpha$  and  $\zeta$  at  $a$ .

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	a	b	z
Step 1	$\alpha, \zeta$	$\beta$	
Step 2	$\alpha$	$\beta$	$\zeta$



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- If broken in favor  $\alpha$ :

	a	b	z
Step 1	$\alpha, \zeta$	$\beta$	
Step 2	$\alpha$	$\beta$	$\zeta$

- If broken in favor  $\zeta$ :

	a	b	z
Step 1	$\zeta, \alpha$	$\beta$	
Step 2	$\zeta$	$\alpha, \beta$	
Step 3	$\beta, \zeta$	$\alpha$	
Step 4	$\beta$	$\alpha$	$\zeta$

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	a	b	z
Step 1	$\alpha, \zeta$	$\beta$	
Step 2	$\alpha$	$\beta$	$\zeta$

Pareto efficient match

- If broken in favor  $\zeta$ :

	a	b	z
Step 1	$\zeta, \alpha$	$\beta$	
Step 2	$\zeta$	$\alpha, \beta$	
Step 3	$\beta, \zeta$	$\alpha$	
Step 4	$\beta$	$\alpha$	$\zeta$

Pareto inefficient match

# Example (Erdil-Ergin): DA-STB

- Outcome of DA-STB is a random matching:

$$\begin{aligned}\mu_{DA-STB} : \alpha &\rightarrow \frac{1}{2}a, \frac{1}{2}b \\ \beta &\rightarrow \frac{1}{2}b, \frac{1}{2}a \\ \zeta &\rightarrow z\end{aligned}$$

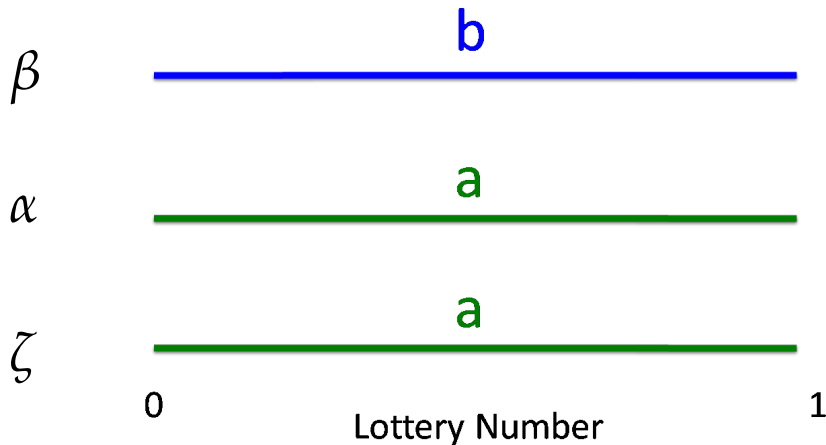
- In the inefficient match  $\alpha, \beta$  form an SIC, i.e. they should trade
- The SOSM outcome (under truth-telling) is:

$$\begin{aligned}\mu_* : \alpha &\rightarrow a \\ \beta &\rightarrow b \\ \zeta &\rightarrow z\end{aligned}$$

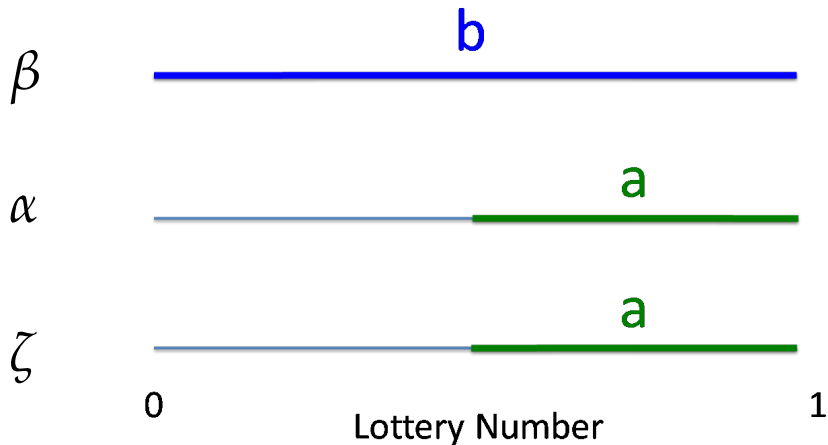
# Example (Erdil-Ergin): Continuum DA-STB

- Unit mass of each type:  $m(\alpha) = m(\beta) = m(\zeta) = 1$
- Break ties by drawing a single lottery number.  
The mass of students has lottery numbers uniformly distributed  
 $\ell \sim U[0, 1]$
- School have capacity  $q_a = q_b = 1, q_z = \infty$

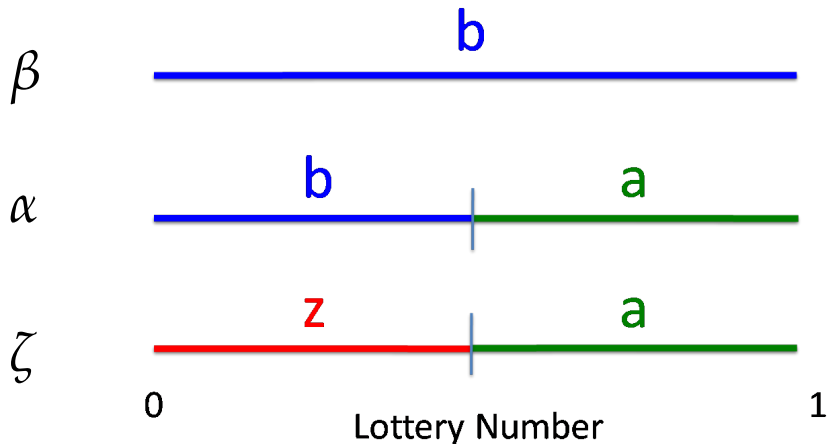
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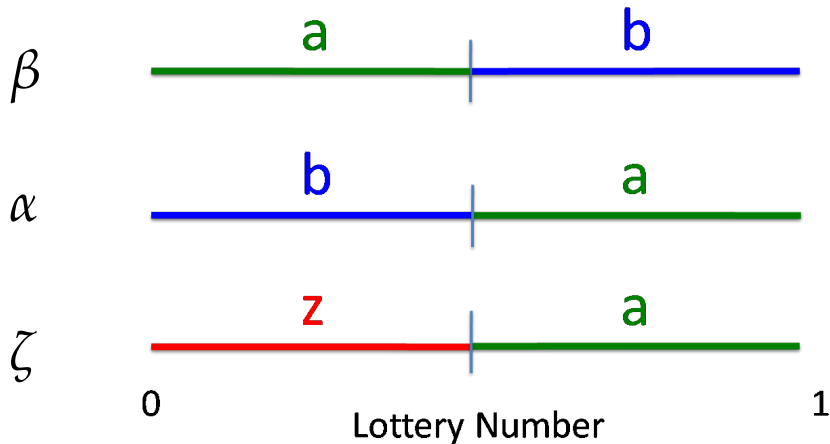
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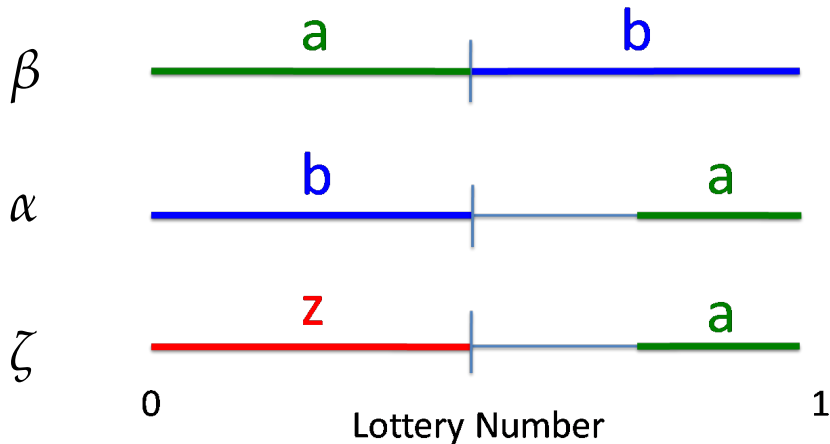


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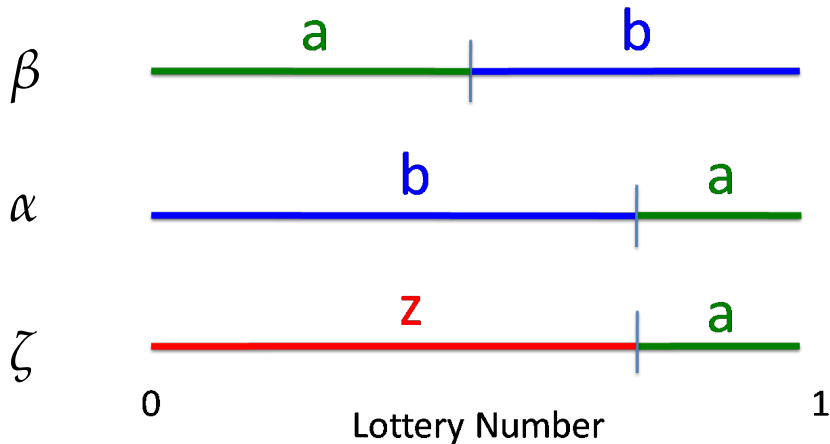




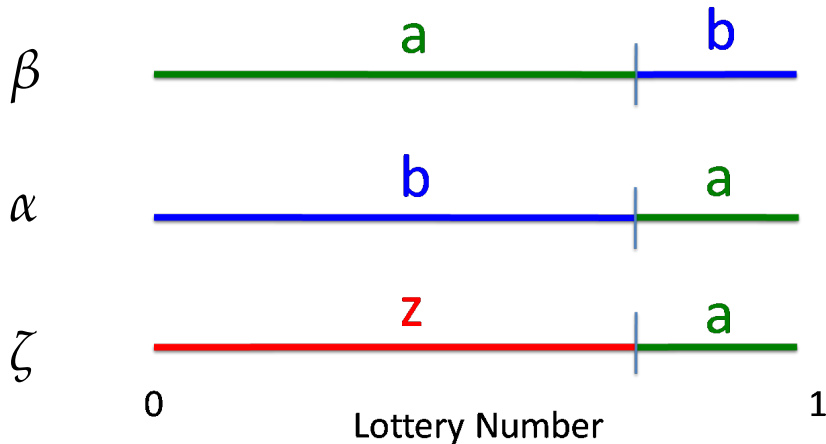
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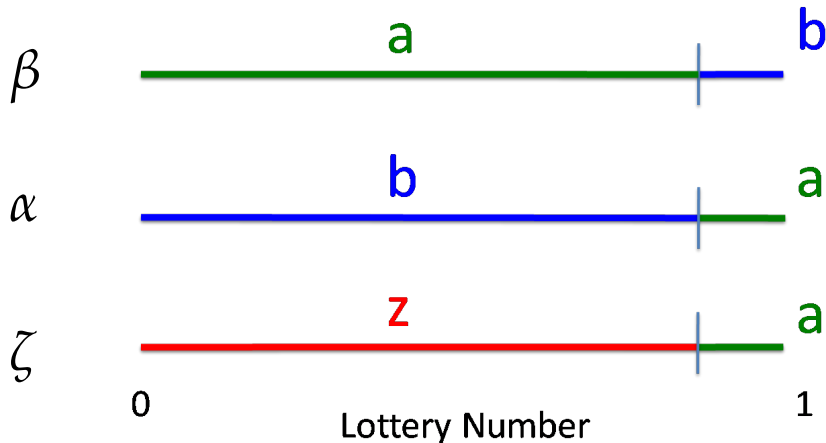
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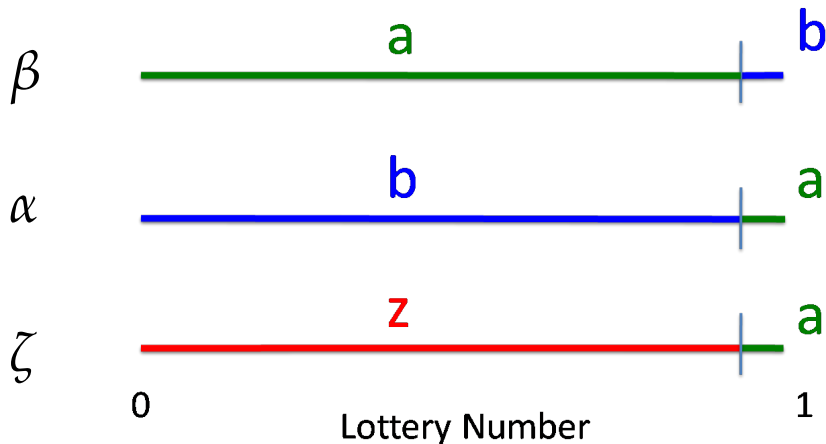
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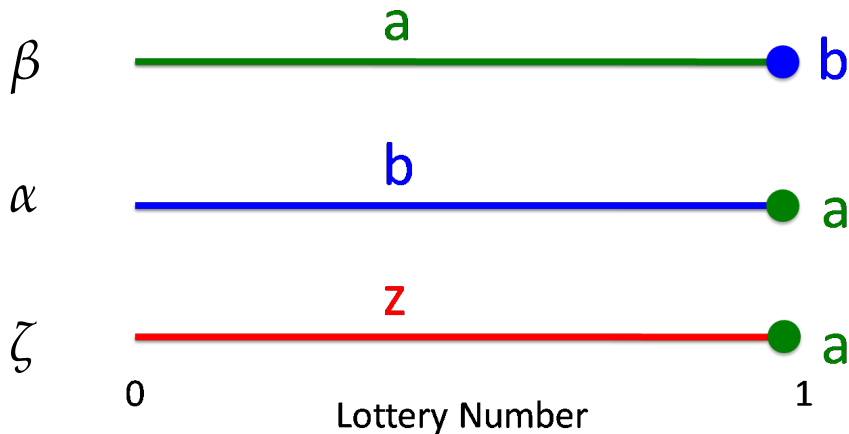
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# Example (Erdil-Ergin): Continuum DA-STB



*cutoffs* are a useful way to represent stable matches

- A system of *cutoffs* is a vector  $p \in \mathbb{R}_+^{|S|}$
- A school  $a$  is *attainable* for student  $\theta$  if  $e_\theta(a) \geq p(a)$
- Given cutoffs  $p$ , student  $\theta$ 's *demand* is

$$D(p, \theta) = \arg \max_{\prec} (\{a \in S \mid e(a) \geq p(a)\} \cup \{\phi\})$$

- The aggregate demand for a school  $a \in S \cup \{\phi\}$  is:

$$D_a(p) = \eta(\{\theta \in \Theta \mid D(p, \theta) = a\})$$

- $p$  is a system of *equilibrium cutoffs* if

$$D_a(p) \leq q_a$$

for all  $a \in S$ , with  $p(a) = 0$  if the inequality is strict

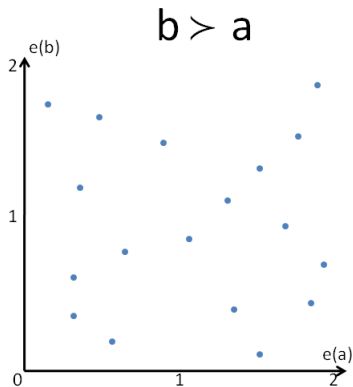
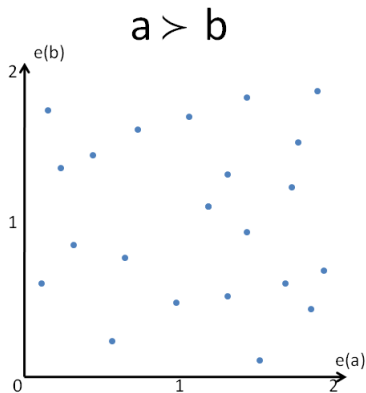
## Lemma ( $\mu \equiv p$ )

*Equilibrium cutoffs are equivalent to stable matching:*

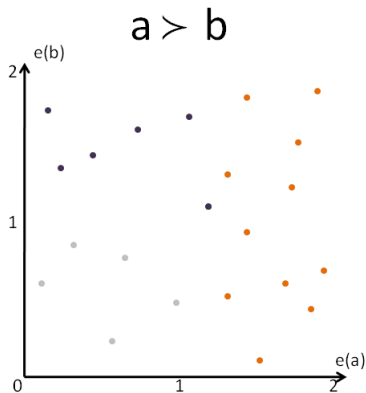
- *The demand of an equilibrium cutoff defines a stable matching.*
- *Every stable matching can be given as the demand of some equilibrium cutoffs.*



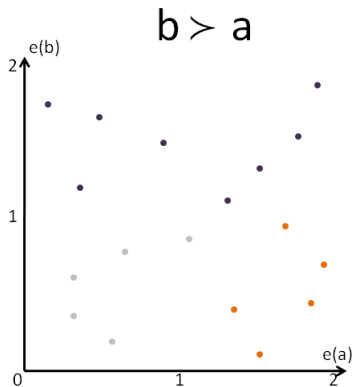
# Cutoffs



# Cutoffs

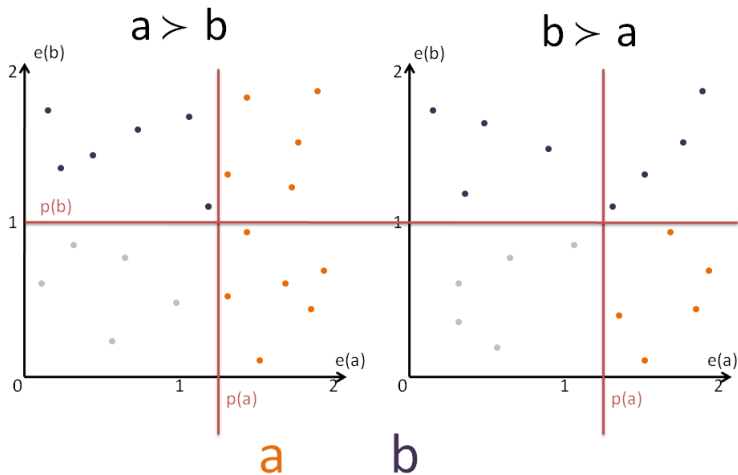


a



b

# Cutoffs



# Using the continuum Framework

- we can describe the stable match using cutoffs  $(p_a, p_b)$ , found by market clearing equations (A-L 2010):

$$\begin{aligned}m(\alpha)(1 - p_a) + m(\beta) \cdot p_b + m(\zeta)(1 - p_a) &= q_a = 1 \\m(\alpha) \cdot p_a + m(\beta)(1 - p_b) &= q_b = 1\end{aligned}$$

- Unique solution:

$$p_a = p_b = 1$$

- The continuum matching:

$$\mu_{DA-STB} : \alpha \rightarrow b, \beta \rightarrow a, \zeta \rightarrow z$$

## Theorem

*When the continuum has a unique stable matching, the DA-STB outcome of a close large matching problem is close to the continuum matching (A-L 2010)*

- Students will misreport a preferences for schools that can be traded
- Over-demanded schools will "enable good trades"
- Congestion, negative externalities
- Students evaluate manipulations ex-ante, considering lotteries over trade.  
Fail to trade ex-post  $\Rightarrow$  inefficiency, instability

# Ex: Pareto Inefficiency

Two special schools:  $a = \text{Art}$ ,  $s = \text{Science}$

Capacities:  $q_a = 1, q_s = 2$

Students:  $m(\tau) = m(\zeta) = 1$  and  $m(\gamma_a) = 2$

School Priorities

0	a	s
$\tau$	1	0
$\zeta$	0	0
$\gamma_a$	0	0
$\gamma_{as}$	0	0

Student Preferences

$\tau$	$\zeta$	$\gamma_a$	$\gamma_{as}$
$s$	$s$	$a$	$a$
$a$	$a$	$\phi$	$s$
$\phi$	$\phi$		$\phi$

$\gamma_a$  is an EU-maximizer with

$$u_{\gamma_a}(a) = 1 > u_{\gamma_a}(\phi) = 0 > u_{\gamma_a}(s) \geq -0.5$$

# Ex: Pareto Inefficiency

## Lemma

*The unique equilibrium outcome of SOSM is Pareto dominated by the DA-STB outcome*



# Ex: Pareto Inefficiency

- $\tau$  and  $\zeta$  report truthfully in equilibrium
- $\gamma_a$  can misreport his type to be  $\gamma_{as}$

Suppose a mass  $v$  of  $\gamma_a$  report  $\gamma_{as}$ :

Type	Preferences	Priority at	Mass
$\tau$	$s \succ a$	$a$	1
$\zeta$	$s \succ a$	—	1
$\gamma_a$	$a$	—	$2 - v$
$\gamma_{as}$	$a \succ s$	—	$v$

## Ex: Pareto Inefficiency

We solve for the DA-STB allocation using the market clearing equations:

$$\begin{aligned}m(\tau) \cdot p_s + m(\zeta)(p_s - p_a)^+ + m(\gamma_{as})(1 - p_a) + m(\gamma_a)(1 - p_a) &= q_a = 1 \\m(\tau)(1 - p_s) + m(\zeta)(1 - p_s) + m(\gamma_{as})(p_a - p_s)^+ &= q_s = 2\end{aligned}$$

The unique solution is:

$$p_a = \frac{v+2}{v+4}, p_s = \frac{v}{v+4}$$

# Ex: Pareto Inefficiency

The DA-STB allocation is:

$$\mu_{DA-STB} : \tau \rightarrow \frac{4}{v+4}s, \frac{v}{v+4}a$$

$$\zeta \rightarrow \frac{4}{v+4}s, \frac{v}{v+4}\phi$$

$$\gamma_a \rightarrow \frac{2}{v+4}a, \frac{v+2}{v+4}\phi$$

$$\gamma_{as} \rightarrow \frac{2}{v+4}a, \frac{2}{v+4}s, \frac{v}{v+4}\phi$$

# Ex: Pareto Inefficiency

The unique Pareto dominating efficient allocation is:

$$\begin{aligned}\mu_* : \tau &\rightarrow s \\ \zeta &\rightarrow \frac{4}{v+4}s, \frac{v}{v+4}\phi \\ \gamma_a &\rightarrow \frac{2}{v+4}a, \frac{v+2}{v+4}\phi \\ \gamma_{as} &\rightarrow \frac{3}{v+4}a, \frac{1}{v+4}s, \frac{v}{v+4}\phi\end{aligned}$$

# Ex: Pareto Inefficiency

The best response of  $\gamma_a$  under SOSM:

- If reports truthfully:

$$\frac{2}{v+4}u_{\gamma_a}(a) + \frac{v+2}{v+4}u_{\gamma_a}(\phi) = \frac{2}{v+4}$$

- If reports  $\gamma_{as}$ :

$$\frac{3}{v+4}u_{\gamma_a}(a) + \frac{1}{v+4}u_{\gamma_a}(s) + \frac{v}{v+4}u_{\gamma_a}(\phi) = \frac{2.5}{v+4}$$

$\Rightarrow$  reporting  $\gamma_{as}$  is a dominant strategy.

# Ex: Pareto Inefficiency

Under DA-STB the unique equilibrium has  $m(\gamma_{as}) = 0$

Under SOSM the unique equilibrium has  $m(\gamma_{as}) = 2$

$\mu_{DA-STB} :$

$$\tau \rightarrow s$$

$$\zeta \rightarrow s$$

$$\gamma_a[\gamma_a] \rightarrow \frac{1}{2}a, \frac{1}{2}\phi$$

$\mu_* :$

$$\tau \rightarrow s$$

$$\zeta \rightarrow \frac{2}{3}s, \frac{1}{3}\phi$$

$$\gamma_a[\gamma_{as}] \rightarrow \frac{1}{2}a, \frac{1}{6}s, \frac{1}{3}\phi$$

# Ex: Pareto Inefficiency - Comments

## Manipulations

- are simple: “grab” a school in order to “trade” it
- work even in a large market, rely only on aggregate information

Students may get “stuck” with a school. Therefore the allocation is

- Pareto inefficient wrt to the true preferences
- unstable and not individually-rational wrt to the true preferences

# The General case

DA-STB is strategy proof  $\Rightarrow$  rank schools by their value

SOSM can be implemented by DA-STB + trade stage

- DA-STB allocation is a lottery of trade opportunities
- The utility of a DA-STB allocation is the utility of the trade lottery, *not* the utility of a final allocation  
 $\Rightarrow$  misreporting
- Students take lotteries with unacceptable schools  
 $\Rightarrow$  positive probability of instability



## Student optimal stable mechanisms:

- ① Can produce unstable outcomes
  - ② Can:
    - Pareto improve upon DA-STB
    - Be Pareto dominated by DA-STB
    - Be not Pareto comparable to DA-STB
  - ③ Susceptible to (simple) manipulations
  - ④ Manipulable in large markets
- ◆ Cutoffs are useful for analyzing school choice problems